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51344 7590 05/25/2010 BROOKS KUSHMAN P.C. /Oracle America/ SUN / STK 1000 TOWN CENTER, TWENTY-SECOND FLOOR SOUTHFIELD, MI 48075-1238			EXAMINER HOANG, HIEU T	
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

DETAILED ACTION

1. This office action is in response to the amendment filed on 04/27/2010.
2. Claims 1-27 are pending.

Response to Arguments

3. Applicant's arguments have been fully considered but they are not persuasive. Applicant argues that the prior art does not teach preempting of data accessing task's resources between a response and a subsidiary task request, and similarly for a subsidiary response and a subsequent subsidiary request, meaning preempting task's resources in pairs of request and response. The examiner respectfully disagrees. Saha in [0059-0060] clearly discloses that upon data being transferred from the sender (a server that data is being read from for example) to the client in buffers (or chunks) and used at the client side, the client acknowledges the transferor that the buffer is now empty (preempting the memory resources associated with the complete data chunk of one task transferred) before new subsequent data chunk can be transferred. Given the broadness of the claimed context resources, the context resources can be read no more than allocated memory blocks in Saha for transferring (DMA) data blocks between a client and a server. The memory blocks are pre-emptible once the data block transfer is finished, or can be emptied to support subsequent tasks or other data blocks of a same task.
4. Applicant argues that the prior art does not teach a channel adapter in the first node for preempting context resources or performing context switch. A channel adapter,

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given broadest reasonable interpretation, can be a module to allocate memory blocks on the client machine, (Saha, [0059-0061], fig. 5, a client executing a read command for a large block of data from the DB server has its memory divided into buffers (memory blocks) fitting a chunk of data; the client will send an ACK to the DB server upon receiving successfully a data chunk and emptying its buffer for receiving another data chunk from the DB server). The claimed channel adapter can be read as than a module (software and/or hardware) on the client to empty the its buffers and send ACKs to the DB server (context switch)

Claim Rejections - 35 USC § 103

5. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

6. Claims 1-11, 13-23, and 25-27 are rejected under 35 U.S.C. 103(a) as being unpatentable over Saha et al. (US 2004/0117375, hereafter Saha), in view of Boyd et al. (US 2004/0049580, hereafter Boyd).

7. For claim 1, Saha discloses a method for performing remote access commands between nodes, the method comprising:

issuing a first request from a first node to a second node, the first request requesting a data access be performed between the first node and the second node ([0060], fig. 5, client writes data to server's database);

receiving, at the first node, a first response from the second node that partially completes the data access ([0060], fig. 5, acknowledgement of write completion associated with a buffer segment at the server);

issuing at least one subsidiary request from the first node to the second node to further complete the data access between the first node and the second node, the at least one subsidiary request based on an amount of partial completion of the data access between the first node and the second node ([0060], fig. 5, after receiving an ACK indicating a buffer read completion from the server, the client can refill its buffer and initiate another subsequent write request); and

receiving, from the second node in response to the at least one subsidiary request, at least one corresponding subsidiary response that further completes the data access between the first node and the second node (fig. 5, another ACK for completion of the subsequent read request);

wherein context resources in a channel adapter of the first node used to service the data access task are pre-emptible after only partial completion of the data access task ((fig. 2, 5, [0059-0061], memory resource preempted for a data chunk when the data chunk has been read or written and acknowledged, the data chunk is the partial task) between the first response and a first subsidiary request for use to support other data access tasks by the channel adapter performing a context switch, and context

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resources in the channel adapter used to service the data access task are pre-emptible between a subsidiary response and a subsequent subsidiary request for use to support other data access tasks by the channel adapter performing a context switch (fig. 2, 5, [0059-0061], memory resource preempted for a data chunk between an ACK response and transferring of the next data block; the client software for emptying its buffers and sending ACKs for context switching is read as the channel adapter).

Saha does not explicitly disclose a data access task.

However, Boyd discloses a data access task (fig. 4, a work request is a task comprising a plurality of data segments or frames)

Therefore, it would have been obvious for one skilled in the art at the time of the invention to combine the teachings of Saha and Boyd to apply the tasking scheme of Boyd to the DMA buffering management scheme of Saha to monitor or manage completion of task segments.

8. Claims 26 and 27 are rejected for the same rationale given in claim 1.

9. For claim 25, Saha discloses a method for performing read remote access commands between nodes ([0022]), the method comprising:

issuing a first read remote access command from a first node to a second node (fig. 5, [0061], read RDMA);

receiving at least one first response from the second node that partially completes a data access associated with the first remote access command ([0061], receiving response with data amount in buffer size);

pre-empting context resources associated with issuance of the first remote access command from the first node prior to completion of the data access associated with the first remote access command (fig. 1, 5, [0059]-[0061], client emptying its allocated buffer blocks after finishing reading the received data);

issuing a second remote access command from the first node using the pre-empted context resources ([0061], client preparing an ACK to the server for each available emptied client buffer); and

issuing, from the first node to the second node, a series of subsidiary remote access commands derived from the first remote access command to receive corresponding subsidiary responses from the second node to complete the data access associated with the first read remote access command ([0061], client sending one or more an ACKs to the server indicating more available client buffers so that more read data can be sent to the client).

Saha does not explicitly disclose a data access task.

However, Boyd discloses a data access task (fig. 4, a work request is a task comprising a plurality of data segments or frames)

Therefore, it would have been obvious for one skilled in the art at the time of the invention to combine the teachings of Saha and Boyd to apply the tasking scheme of

Boyd to the DMA buffering management scheme of Saha to monitor or manage completion of task segments.

10. For claim 2, Saha-Boyd discloses the invention as in claim 1. Saha-Boyd further discloses issuing the first request from the first node comprises:

detecting an application request in a request queue, the application request identifying the data access task to be performed between the first and second node (RDMA read/write request); and wherein the method comprises: repeating issuing at least one subsidiary request and receiving at least one corresponding subsidiary response between the first and second nodes until the data access task is totally complete between the first and second nodes (Saha, [0063], queue segments for subsequent buffer read/write iteration).

11. For claim 3, Saha-Boyd discloses the invention as in claim 2. Saha-Boyd further discloses allocating context resources in the first node for receipt of the first response and for receipt of the at least one subsidiary response ([0061], client buffers for receiving read data responses), the context resources allocated to support receipt of the first response and for receipt of the at least one subsidiary response that contain data in an amount not exceeding a data allotment to support at least partial completion of the data access task (Saha, fig. 5, [0060], client buffers of fix size for receiving read data segments or a partial task).

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12. For claim 4, Saha-Boyd discloses the invention as in claim 3. Saha-Boyd further discloses:

pre-empting the context resources in the first node allocated for receipt of the first response and for receipt of at least one subsidiary response prior to full completion of the data access task; issuing a second request from the first node by re-activating the context resources that were pre-empted; and (Saha, [0060], client completes reading from a buffer and sends ACK indication that buffer is available for more data);

at a time after issuance of the second request from the first node, allocating context resources to support at least partial completion of the data access task in order to continue repeating issuing at least one subsidiary request and receiving at least one corresponding subsidiary response between the first and second nodes until the data access task is totally complete between the first and second nodes (Saha, [0060], after a partial read completion, send ACK from client to server indicating more data can be sent to the available buffer)

13. For claim 5, Saha-Boyd discloses the invention as in claim 4. Saha-Boyd further discloses allocating context resources to support at least partial completion of the data access task comprises:

reserving the context resources to process subsidiary responses containing data limited in amount according to a preset data allotment identifying an amount of data to be transferred between the first and second nodes (Saha, fig. 5, buffer of fixed size).

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14. For claim 6, Saha-Boyd discloses the invention as in claim 1. Saha-Boyd further discloses the second request is at least one of: a different type of data access request than the first request; and received in a request queue that is different than a request queue from which the first request was received (Saha, [0063], [0060], [0061], subsequent queued requests can be read or write).

15. For claim 7, Saha-Boyd discloses the invention as in claim 4. Saha-Boyd further discloses issuing at least one subsidiary request comprises:

calculating a remaining amount of data required to complete the data access task between the first node and the second node (Saha, [0063], queueing remaining data or task segments to be executed, Boyd task consisting of many segments); and

creating at least one subsidiary request to reference at least a portion of the remaining amount of data required to complete the data access task (Saha, [0063], [0060], [0061], subsequent requests amount is the amount of empty buffers)

16. For claim 8, Saha-Boyd discloses the invention as in claim 7. Saha-Boyd further discloses calculating the remaining amount of data comprises:

determining a total completed amount of data processed for the data access task by the first request and associated first response and all subsidiary requests and corresponding subsidiary responses between the first and second node (Saha, [0061], initial total amount of data to be executed in subsequent buffer transfers in a queue); and

determining the remaining amount of data required to complete the data access task as a difference between an initial amount of data specified by an application request and the total completed amount of data (Saha, [0061], remaining amount of data of a task is the remaining data in the queue and is the difference of initial amount and complete amount).

17. For claim 9, Saha-Boyd discloses the invention as in claim 4. Saha-Boyd further discloses:

the first and second nodes are nodes that utilize channel adapters to exchange the first request and the at least one subsidiary request and the corresponding first response and the at least one subsidiary response (Saha, [0022], Infiniband);

the application request is a remote direct memory access request for the first node to access data in a memory at the second node (Saha, [0022], Remote DMA); and

an initial amount of data specified by the application request is a total amount of data that the first node is to access in the memory at the second node (Saha, a RDMA request inherently has a task size to put in a queue).

18. For claim 10, Saha-Boyd discloses the invention as in claim 9. Saha-Boyd further discloses the first request and the at least one subsidiary request are read remote direct memory access commands issued by the first node to read data in the memory from the second node ([0061], RDMA read).

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19. For claim 11, Saha-Boyd discloses the invention as in claim 1. Saha-Boyd further discloses:

establishing a data allotment as a maximum amount of data to be used when responding to requests to transfer portions of data between the nodes, such that if a total amount of data to be transferred between the first node and the second node is greater than the data allotment, the second node provides the first response and the at least one subsidiary response that contain response data that does not exceed the data allotment (Saha, fig. 5, response in fixed size allotments or window); and

allocating context resources in the first node for receipt of the first response and for receipt of the at least one subsidiary response, the context resources allocated to support receipt of responses in an amount that does not exceed the data allotment (Saha, [0061], receive read data in amounts no more than a buffer size).

20. Claims 13-23 are rejected for the same rationale given in claims 1-11 respectively.

21. Claims 12 and 24 are rejected under 35 U.S.C. 103(a) as being unpatentable over Saha-Boyd, further in view of Turner et al. (An Approach For Congestion Control in Infiniband, hereafter Turner).

22. For claim 12, Saha-Boyd discloses the invention as in claim 11. Saha-Boyd does not disclose establishing the data allotment comprises dynamically determining the data

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allotment between the first and second nodes based on at least one external data allotment event, such that if the at least one external data allotment event occurs, the first and second nodes change a value of the data allotment.

However, Turner discloses the same (fig. 2, section 5.1.3.2, changing a receive window based on a condition)

Therefore, it would have been obvious for one skilled in the art at the time of the invention to combine the teachings of Saha-Boyd and Turner to adjust the amount of maximum data allotment (window control) to, e.g., avoid network congestion (Turner, abstract).

23. Claim 24 is rejected for the same rationale given in claim 12.

Conclusion

24. **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the mailing date of this final action.

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25. Any inquiry concerning this communication or earlier communications from the examiner should be directed to Hieu T. Hoang whose telephone number is 571-270-1253. The examiner can normally be reached on Monday-Thursday, 8 a.m.-5 p.m., EST.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Thu Nguyen can be reached on 571-272-6967. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300. Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/HH/

Patent Examiner, Art Unit 2452

/DOHM CHANKONG/

Primary Examiner, Art Unit 2452